

Beyond Johnson–Lindenstrauss

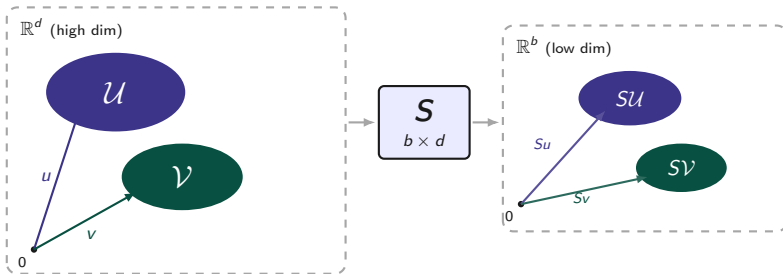
Uniform Bounds for Sketched Bilinear Forms

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Bound how much a sketch distorts inner products



Goal

Bound, with high probability,

$$\sup_{u \in \mathcal{U}, v \in \mathcal{V}} |\langle Su, Sv \rangle - \langle u, v \rangle|.$$

— motivation —

Where this template shows up

$$\text{Template: } \sup_{u \in \mathcal{U}, v \in \mathcal{V}} |u^\top S^\top S v - u^\top v|.$$

classical — $u = v$

bilinear — $u \neq v$

Johnson–Lindenstrauss lemma

Preserve norms on a finite set $\mathcal{A}$:

$$| \|Su\|^2 - \|u\|^2 |, \quad u \in \mathcal{A}.$$

$$\mathcal{U} = \mathcal{V} = \mathcal{A}$$

Restricted isometry property

Same, but over s -sparse vectors Σ_s :

$$| \|Su\|^2 - \|u\|^2 |, \quad u \in \Sigma_s.$$

$$\mathcal{U} = \mathcal{V} = \Sigma_s$$

Linear regression

Model $y_i = \beta^\top x_i + \eta_i$, $\beta \in \mathcal{B}$, $x_i \in \mathcal{X}$.

Bound $|\beta^\top S^\top S x - \beta^\top x|$.

Federated learning

Sketch gradient $g \in \mathcal{G}$; Hessian eigvec $h \in \mathcal{H}$.

Bound $|g^\top S^\top S h - g^\top h|$.

Linear bandits

Reward $r_i = \theta_\star^\top a_i + \eta_i$, $\theta_\star \in \Theta$, $a_i \in \mathcal{A}$.

Bound $|\theta_\star^\top S^\top S a - \theta_\star^\top a|$.

Random quadratic forms: from fixed A to a set $\mathcal{A}$

$$\sup_{u,v} |\langle Su, Sv \rangle - \langle u, v \rangle| \iff C_{\mathcal{A}}(\xi) = \sup_{A \in \mathcal{A}} |\|A\xi\|_2^2 - \mathbb{E}\|A\xi\|_2^2| = \sup_{A \in \mathcal{A}} |\xi^\top A^\top A \xi - \mathbb{E} \xi^\top A^\top A \xi|.$$

Hanson–Wright (fixed matrix A)

For sub-Gaussian ξ with i.i.d. entries,

$$\mathbb{P}(|\|A\xi\|_2^2 - \mathbb{E}\|A\xi\|_2^2| \geq \epsilon) \lesssim \exp\left(-c \min\left(\frac{\epsilon^2}{\|A^\top A\|_F^2}, \frac{\epsilon}{\|A^\top A\|_{2 \rightarrow 2}}\right)\right).$$

Krahmer–Mendelson–Rauhut '14 (uniform over a set $\mathcal{A}$)

Let $C_{\mathcal{A}}(\xi) = \sup_{A \in \mathcal{A}} |\|A\xi\|_2^2 - \mathbb{E}\|A\xi\|_2^2|$. Define

$$W = \gamma_2(\mathcal{A})(\gamma_2(\mathcal{A}) + d_F(\mathcal{A})) + d_F(\mathcal{A})d_{2 \rightarrow 2}(\mathcal{A}), \quad V = d_{2 \rightarrow 2}(\mathcal{A})(\gamma_2(\mathcal{A}) + d_F(\mathcal{A})), \quad U = d_{2 \rightarrow 2}^2(\mathcal{A}).$$

$$\mathbb{P}(C_{\mathcal{A}}(\xi) \geq c_1 W + \epsilon) \leq 2 \exp\left(-c_2 \min\left(\frac{\epsilon^2}{V^2}, \frac{\epsilon}{U}\right)\right).$$

— main result —

A bilinear Hanson–Wright for double sets

Theorem 3.1

Let $\mathcal{M} \subset \mathbb{R}^{m \times n}$, $\mathcal{N} \subset \mathbb{R}^{n \times m}$, ξ sub-Gaussian with i.i.d. entries, and

$$C_{\mathcal{M}, \mathcal{N}}(\xi) = \sup_{M \in \mathcal{M}, N \in \mathcal{N}} \left| \xi^\top M^\top N \xi - \mathbb{E} \xi^\top M^\top N \xi \right|.$$

$$W = \gamma_2(\mathcal{M}) \cdot [\gamma_2(\mathcal{N}) + d_F(\mathcal{N})] + \gamma_2(\mathcal{N}) \cdot [\gamma_2(\mathcal{M}) + d_F(\mathcal{M})],$$

$$V = \gamma_2(\mathcal{M}) d_{2 \rightarrow 2}(\mathcal{N}) + \gamma_2(\mathcal{N}) d_{2 \rightarrow 2}(\mathcal{M}) + \min\{d_F(\mathcal{M}) d_{2 \rightarrow 2}(\mathcal{N}), d_F(\mathcal{N}) d_{2 \rightarrow 2}(\mathcal{M})\},$$

$$U = d_{2 \rightarrow 2}(\mathcal{M}) d_{2 \rightarrow 2}(\mathcal{N}).$$

$$\mathbb{P}(C_{\mathcal{M}, \mathcal{N}}(\xi) \geq c_1 W + \epsilon) \leq 2 \exp\left(-c_2 \min\left(\frac{\epsilon^2}{V^2}, \frac{\epsilon}{U}\right)\right).$$

Take-aways. (i) Setting $\mathcal{M} = \mathcal{N}$ recovers Krahmer et al. (ii) Complexities of $\mathcal{M}, \mathcal{N}$ enter *separately* — no joint set $\{M^\top N\}$. (iii) If one set is small, the bound is small.

— proof sketch —

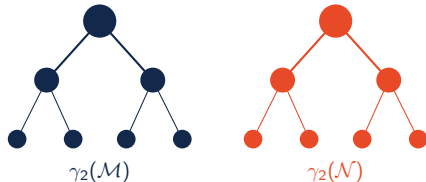
Key technique: *double chaining*

Chain twice. Two admissible sequences $\{T_r(\mathcal{M})\}$, $\{T_r(\mathcal{N})\}$. Each chain pairs its increments with the worst case of the other set.

Double-tree lemma (Thm. 3.2)

If $\{X_{(u,v)}\}$ is sub-Gaussian-Lipschitz in each coordinate,

$$\sup_{u \in \mathcal{U}, v \in \mathcal{V}} |X_{(u,v)}| \lesssim C_u \gamma_2(\mathcal{U}, d) + C_v \gamma_2(\mathcal{V}, d).$$



$$\text{bound} \leq C \cdot (\gamma_2(\mathcal{M}) + \gamma_2(\mathcal{N}))$$

— applications —

What the new bound buys us

01 Johnson–Lindenstrauss

Classical lemma drops out as a special case of our bilinear bound.

02 RIP for inner products

For arbitrary $\mathcal{U}, \mathcal{V}$ on the sphere, RIP holds for

$$n \gtrsim \frac{1}{\epsilon^2} (\omega(\mathcal{U}) + \omega(\mathcal{V}))^2.$$

Gaussian-width rate; matches $\sqrt{\log |\mathcal{T}|}$ on finite sets.

03 Sketched federated learning

Reuse a single sketching matrix across rounds. Convergence depends on gaussian width of gradient & loss-Hessian eigenvector sets.

04 Sketched linear regression

Solve Least Squares in the sketched space. Error bound depends on $\omega(\mathcal{B})$ and $\omega(\mathcal{X})$.

— extensions & experiments —

Beyond i.i.d., and a sketched bandit algorithm

Conditionally independent coordinates

Same bounds hold under a relaxed assumption that covers **CountSketch** and **SRHT** (Subsampled Randomized Hadamard Transform).

Sum of T random quadratic forms

$$\sup_{M, N} \left| \sum_{t=1}^T \xi_t^\top M^\top N \xi_t - \mathbb{E}[\xi_t^\top M^\top N \xi_t] \right|$$

Scales with extra $\sqrt{T}$. Needed when the sketching matrix changes per round.

Sketched linear bandits

We design compressed versions of linear bandit algorithms, where actions are projected to a lower-dimensional sketched space before applying LinUCB-style exploration.

